

Today's outline - April 19, 2022



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- Mixtures of correctable errors

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Homework Assignment #07:
Chapter 9:2,3,4; Chapter 10:3,4,11
Due Thursday, April 21, 2022



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Homework Assignment #07:
Chapter 9:2,3,4; Chapter 10:3,4,11
Due Thursday, April 21, 2022

Exam #2 - April 26, 2022



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- NMR-based quantum computers
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Final Exam will be 10 minute presentations
on a quantum computing journal article

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Measurement with O will give the state $P_i \rho' P_i^\dagger$ with probability $|\alpha_i|^2$ and will leave the system in the pure state $E_i|c\rangle$

$$P_i \rho' P_i^\dagger = E_i |c\rangle\langle c| E_i^\dagger$$

Knowing the result of the measurement, λ_i , means that the error can be corrected by application of E_i^\dagger thus recovering the corrected state $|c\rangle$

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$$\langle c_1 | X_i^\dagger Z_i | c_2 \rangle = \langle c_1 | Z_i^\dagger X_i | c_2 \rangle = \langle c_1 | I Z_i | c_2 \rangle = \langle c_1 | I X_i | c_2 \rangle = 0$$

In addition, it is possible to show that Y_i are also correctable errors using multiplication in the Pauli group and the commutator relations

$$[X, Y] = XY - YX = -2Z, \quad [Y, Z] = YZ - ZY = -2X, \quad [Z, X] = ZX - XZ = 2Y$$

$$X_i^\dagger Y_i = X_i^\dagger (-X_i Z_i) = -I Z_i$$

$$\langle c_1 | X_i^\dagger Y_i | c_2 \rangle = -\langle c_1 | I Z_i | c_2 \rangle = 0$$

$$Z_i^\dagger Y_i = Z_i^\dagger (Z_i X_i) = I X_i$$

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So any code that corrects all bit-flip (X) and phase-flip (Z) errors also will correct all Y errors

Requirements for quantum computers



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These hold for the standard circuit model, alternative models require more general criteria

NMR-based quantum computers



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NMR-based quantum computers

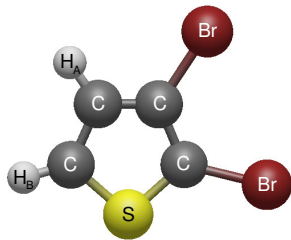


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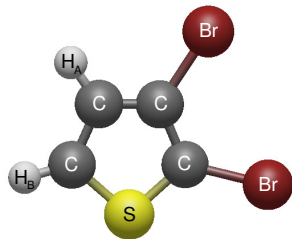
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In a 4.7 T magnetic field, the proton spin states are separated by 200 MHz but are slightly different because of the differing local environment



2-qubit NMR computer



Assume that the state of the entire system can be described as an ensemble of non-interacting molecules (e.g. liquid)

“Bulk spin-resonance quantum computation,” N.A. Gershenfeld and I.L. Chuang, *Science* **275**, 350-356 (1997).

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The **deviation from equilibrium** is what is being measured as the qubit state

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The dynamics of the ensemble can be approximated by just the deviation density matrix whose macroscopic signal has a relatively long decoherence time

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2-qubit NMR computer



At equilibrium, suppose that the deviation density matrix is given by

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At equilibrium, suppose that the deviation density matrix is given by

$$\rho_{\Delta} = \alpha \begin{pmatrix} 3 & 0 & 0 & 0 \\ 0 & -1 & 0 & 0 \\ 0 & 0 & -1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix}$$

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The C_{not} gate can be implemented due to the nonlinear interaction between spins on the same molecule

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5-qubit NMR computer

In 2000, this was extended to a 5-qubit system using the same principles

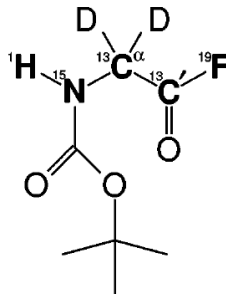
“Approaching five-bit NMR quantum computing,” R. Marx, A.F. Fahmy, J.M. Myers, W. Bermel, and S.J. Glaser, *Phys. Rev A* **62**, 012310 (2000).

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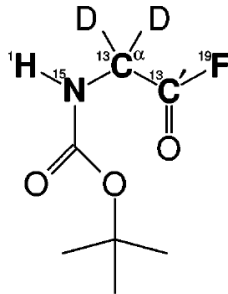


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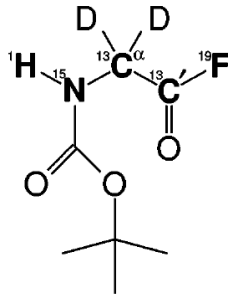
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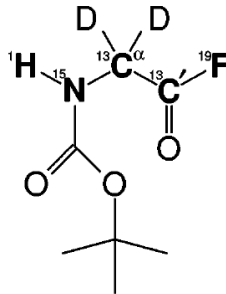


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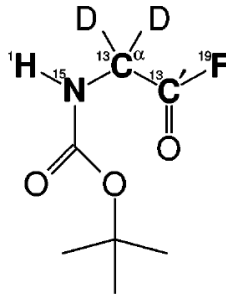
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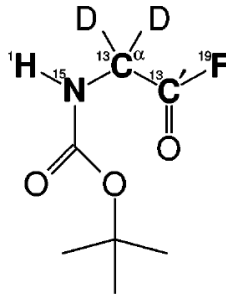


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The authors claim that this computer can demonstrate at least 10 quantum algorithms

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